Frobenius structures

A partial solution of the conjecture 000

Questioning pseudoaffine

Prenuclear vs. nuclear objects in *-autonomous categories

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Frobenius quantales

Definition

A **Frobenius quantale** is a quantale (Q, \star) coming with antitone "negations" $^{\perp}(-), (-)^{\perp} : Q \longrightarrow Q$ satisfying:

> $({}^{\perp}x)^{\perp} = {}^{\perp}(x^{\perp}) = x$ (inverse to each other) $x \multimap {}^{\perp}y = x^{\perp} \multimap y$ (contraposition law)

Remarks

- A provability model of non-commutative classical linear logic.
- Similar to involutive residuated lattice, but complete.
- Such a quantale is until if and only if it has a dualizing element.
- The above axiomatization allows to consider Frobenius quantales with neither a unit nor a dualizing element.

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Motivations

Theorem (Egger, Kruml, Paseka ~ 2008, Santocanale 2020) Let L be a complete lattice. The following are equivalent:

- **1.** The quantale $[L, L]_{\vee}$ of join-preserving endomaps of L is Frobenius.
- 2. L is a completely distributive lattice.

Theorem (Raney 1960, Higgs and Rowe 1989)

The nuclear objects in the category of complete sup-lattices are exactly the completely distributive lattice.

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Motivations

Conjecture

Let A be an object of a SMCC. The following are equivalent:

- 1. The object [A, A] of endomorphisms of A has a Frobenius structure.
- 2. A is nuclear.

Theorem (Raney 1960, Higgs and Rowe 1989)

The nuclear objects of the category of complete sup-lattices are exactly the completely distributive lattices.

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Dual pairings

Definition

A triple (A, B, ϵ) is said to be a *dual pairing* (w.r.t. the object 0) if

 $\epsilon: A \otimes B \longrightarrow 0$

and the two induced natural transformations are isomorphims.

 $\operatorname{hom}(X,B) \longrightarrow \operatorname{hom}(A \otimes X,0), \quad \operatorname{hom}(X,A) \longrightarrow \operatorname{hom}(X \otimes B,0).$

Example

In a *-autonomous category

 $ev: A \otimes A^* \longrightarrow 0, \qquad \epsilon: (A \otimes A^*) \otimes [A, A] \longrightarrow 0$

are dual pairings.

In SLatt, the map

$$\epsilon(x,y) = \begin{cases} \bot, & x \le y, \\ \top, & x \nleq y, \end{cases}$$

yields a dual pairing $\epsilon_L : L \otimes L^{op} \longrightarrow 2$. $(\Box \rightarrow \langle \Box \rangle \land \exists \land \exists \land \exists \land \neg \land \bigcirc$

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• In SLatt, the map

$$\epsilon(x,y) = \begin{cases} \bot, & x \leq y, \\ \top, & x \notin y, \end{cases}$$

yields a dual pairing $\epsilon_L : L \otimes L^{op} \longrightarrow 2$.

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Frobenius quantales, once more

In a Frobenius quantale $(Q, \star, {}^{\perp}(-), (-)^{\perp})$, we have

- (Q, Q^{op}, ϵ) is a dual pairing;
- (Q, *) is a semigroup;
- $^{\perp}(-), (-)^{\perp} : Q \rightarrow Q^{\text{op}} \text{ and } x \leq ^{\perp}y \text{ iff } y \leq x^{\perp};$

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Frobenius structures

Definition A *Frobenius structure* is a tuple $(A, B, \epsilon, \mu_A, l, r)$ where

- (A, B, ϵ) is a dual pairing;
- (A, μ_A) is a semigroup;
- $l, r : A \longrightarrow B$ and (l, r) is an adjunction with both maps invertible

such that

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$$\begin{array}{ccc} A \otimes A & \xrightarrow{A \otimes r} & A \otimes B \\ \downarrow_{\otimes A} & & \downarrow_{\alpha^{\ell}} \\ B \otimes A & \xrightarrow{\alpha^{\rho}} & B. \end{array}$$

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From here, \mathcal{V} is symmetric monoidal closed and 0 = I.

Definition

For every object A of C, there exists a canonical arrow

$$\min_{A} : A^* \otimes A \longrightarrow [A, A].$$

An object A is *nuclear* if mix_A is an isomorphism.

Theorem (Raney 1960, Higgs and Rowe 1989)

The nuclear objects of SLatt are exactly the completely distributive lattices.

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From Nuclearity to Frobenius structure, and back

Theorem (LS and CL, CSL 2023)

If an object A of \mathcal{V} is nuclear, then [A, A] can be endowed with a Frobenius structure.

If \mathcal{V} is *-autonomous and A is pseudoaffine, then the converse hold.

Definition

An object *A* of \mathcal{V} , is **pseudoaffine** if (either is initial or) *I* embeds into *A* as a retract (*i.e* if there exists $p : I \rightarrow A$ and $c : A \rightarrow I$ such that $c \circ p = id_{I}$.).

Example

Every object of SLatt, k-Vect, Coh HypCoh, is pseudoaffine.

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Is pseudoaffine necessary ?

Question:

Can we have a $\ast\text{-autonomous}$ category $\mathcal V$ and a object A of $\mathcal V$ such that

- A is not nuclear,
- A is not pseudoaffine,
- A is prenuclear, that is [A, A] is Frobenius ?

Definition (Schalk and de Paiva 2004 category P-Set)

Let (P, \leq) be a poset (the base category). We define the category *P*-Set:

- An object: a pair (X, α) with X a set and $\alpha : X \to P$;
- An arrow $(X, \alpha) \rightarrow (Y, \beta)$: a relation $R \in P(X \times Y)$ such that xRy implies $\alpha(x) \le \beta(y)$.

Theorem (Schalk and de Paiva 2004)

For $(Q, \star, 1)$ a unital commutative Frobenius quantale, Q-Set is *-autonomous.

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Various characterisations

We take Q such that 1 = 0 in Q, so I = 0 in Q-Set.

Lemmas

A Q-Set (X, α) is

- pseudoaffine iff $\alpha(x) = 1$, for some $x \in X$,
- nuclear iff,

$$\alpha(x) \multimap \alpha(y) = \alpha(x)^{\perp} \star \alpha(y) \tag{(\forall x, y)}$$

Equivalently: iff $\alpha(x)$ is invertible $(\forall x)$

prenuclear iff, for a pair of inverse maps (f, g) on X,

$$\alpha(x) \multimap \alpha(y) = \alpha(fx)^{\perp} \star \alpha(y) = \alpha(x)^{\perp} \star \alpha(gy) \qquad (\forall x, y)$$

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Finding the right Q

- Take ℤ.
- Add ±∞.
- Add a new unit u such that 0 < u < 1.

Then $\mathbb{Z} \subseteq Q$ is prenuclear, with f(x) = x - 1 and g(x) - x + 1.

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Discussing the pseudoaffine condition

We take Q such that 1 = 0 in Q, so I = 0 in Q-Set.

Theorem (LS and CL)

If (X, α) is prenuclear and any of the following conditions holds:

- the quantale Q has no infinite chain,
- X is finite,
- $\alpha(x)$ is invertible, for some $x \in X$,
- (X, α) is Girard, that is, f = g,

then (X, α) is nuclear.

Theorem (LS and CL)

There is a quantale Q and an object (X, α) of Q-**Set** which is prenuclear but not nuclear.

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Conclusions

Results

- A definition of Frobenius structures in autonomous categories.
- Proof of our conjecture up to a technical (but quite natural) hypothesis (pseudoaffine).
- Other results, e.g. an abstraction of the double negation construction.
- Testing the pseudoaffine condition.

To do next

- What about Girard quantales/structures and nuclearity ?
- Understand "how much" we need *-autonomous categories.
- Monoidal fibrations.

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Thank you!

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